

Type theory in type theory using single substitutions

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What is type theory?

What is type theory?

Con : Set

Ty : Con → Set

Var : (Γ : Con) → Ty Γ → Set

Tm : (Γ : Con) → Ty Γ → Set

What is type theory?

$\text{Con} : \text{Set}$

$\text{Ty} : \text{Con} \rightarrow \text{Set}$

$\diamond : \text{Con}$

$\text{Var} : (\Gamma : \text{Con}) \rightarrow \text{Ty } \Gamma \rightarrow \text{Set}$

$\text{Tm} : (\Gamma : \text{Con}) \rightarrow \text{Ty } \Gamma \rightarrow \text{Set}$

$- \triangleright - : (\Gamma : \text{Con}) \rightarrow \text{Ty } \Gamma \rightarrow \text{Con}$

What is type theory?

Con : Set

Ty : Con → Set

◊ : Con

Sub : Con → Con → Set

- [-] : Ty Γ → Sub Δ Γ → Ty Δ

p : Sub (Γ ▷ A) Γ

Var : (Γ : Con) → Ty Γ → Set

Tm : (Γ : Con) → Ty Γ → Set

- ▷ - : (Γ : Con) → Ty Γ → Con

var : Var Γ A → Tm Γ A

vz : Var (Γ ▷ A) (A [p])

vs : Var Γ A → Var (Γ ▷ B) (A [p])

What is type theory?

Con : Set

Ty : Con → Set

◊ : Con

Sub : Con → Con → Set

-[-] : Ty Γ → Sub Δ Γ → Ty Δ

p : Sub (Γ ▷ A) Γ

Π : (A : Ty Γ) → Ty (Γ ▷ A) → Ty Γ

-↑ : (σ : Sub Δ Γ) →

Sub (Δ ▷ A [σ]) (Γ ▷ A)

Var : (Γ : Con) → Ty Γ → Set

Tm : (Γ : Con) → Ty Γ → Set

- ▷ - : (Γ : Con) → Ty Γ → Con

var : Var Γ A → Tm Γ A

vz : Var (Γ ▷ A) (A [p])

vs : Var Γ A → Var (Γ ▷ B) (A [p])

Π[] : (Π A B)[σ] = Π (A [σ]) (B [σ↑])

What is type theory?

Con : Set

Ty : Con → Set

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-[-] : Ty Γ → Sub Δ Γ → Ty Δ

p : Sub (Γ ▷ A) Γ

Π : (A : Ty Γ) → Ty (Γ ▷ A) → Ty Γ

-↑ : (σ : Sub Δ Γ) →

Sub (Δ ▷ A[σ]) (Γ ▷ A)

{-} : Tm Γ A → Sub Γ (Γ ▷ A)

Var : (Γ : Con) → Ty Γ → Set

Tm : (Γ : Con) → Ty Γ → Set

- ▷ - : (Γ : Con) → Ty Γ → Con

var : Var Γ A → Tm Γ A

vz : Var (Γ ▷ A) (A[p])

vs : Var Γ A → Var (Γ ▷ B) (A[p])

Π[] : (Π A B)[σ] = Π (A[σ]) (B[σ↑])

lam : Tm (Γ ▷ A) B → Tm Γ (Π A B)

app : Tm Γ (Π A B) →

(u : Tm Γ A) → Tm Γ (B[⟨u⟩])

What is type theory? cont.

$$\begin{array}{l} -[-] : \text{Tm } \Gamma A \rightarrow \\ (\sigma : \text{Sub } \Delta \Gamma) \rightarrow \text{Tm } \Delta (A[\sigma]) \end{array} \quad \text{app (lam } t) u = t[\langle u \rangle]$$

What is type theory? cont.

$$[-] : \text{Tm } \Gamma A \rightarrow$$

$$(\sigma : \text{Sub } \Delta \Gamma) \rightarrow \text{Tm } \Delta (A[\sigma])$$

$$A[\langle u \rangle][\sigma] = A[\sigma \uparrow][\langle u[\sigma] \rangle]$$

$$\text{app} (\text{lam } t) u = t[\langle u \rangle]$$

$$\underbrace{(\text{lam } t)[\sigma]}_{\text{Tm } \Delta ((\Pi A B)[\sigma])} = \underbrace{\text{lam } (t[\sigma \uparrow])}_{\text{Tm } \Delta (\Pi (A[\sigma]) (B[\sigma \uparrow]))}$$

$$\underbrace{(\text{app } t u)[\sigma]}_{\text{Tm } \Delta (B[\langle u \rangle][\sigma])} = \underbrace{\text{app } (t[\sigma]) (u[\sigma])}_{\text{Tm } \Delta (B[\sigma \uparrow][\langle u[\sigma] \rangle])}$$

What is type theory? cont.

$$-[-] : \text{Tm } \Gamma A \rightarrow$$

$$(\sigma : \text{Sub } \Delta \Gamma) \rightarrow \text{Tm } \Delta (A[\sigma])$$

$$A[\langle u \rangle][\sigma] = A[\sigma \uparrow][\langle u[\sigma] \rangle]$$

$$A[p][\langle u \rangle] = A$$

$$A[p][\gamma \uparrow] = A[\gamma][p]$$

$$\text{app } (\text{lam } t) u = t[\langle u \rangle]$$

$$\frac{(\text{lam } t)[\sigma] = \text{lam } (t[\sigma \uparrow])}{\text{Tm } \Delta ((\Pi A B)[\sigma]) \quad \text{Tm } \Delta (\Pi (A[\sigma]) (B[\sigma \uparrow]))}$$

$$\frac{(\text{app } t u)[\sigma] = \text{app } (t[\sigma])(u[\sigma])}{\text{Tm } \Delta (B[\langle u \rangle][\sigma]) \quad \text{Tm } \Delta (B[\sigma \uparrow][\langle u[\sigma] \rangle])}$$

$$(\text{var } x)[p] = \text{var } (\text{vs } x)$$

$$(\text{var } \text{vz})[\langle u \rangle] = u$$

$$(\text{var } (\text{vs } x))[\langle u \rangle] = \text{var } x$$

$$(\text{var } \text{vz})[\gamma \uparrow] = \text{var } \text{vz}$$

$$(\text{var } (\text{vs } x))[\gamma \uparrow] = (\text{var } x)[\gamma][p]$$

What is type theory? cont.

$$[-] : \text{Tm } \Gamma A \rightarrow$$

$$(\sigma : \text{Sub } \Delta \Gamma) \rightarrow \text{Tm } \Delta (A[\sigma])$$

$$A[\langle u \rangle][\sigma] = A[\sigma \uparrow][\langle u[\sigma] \rangle]$$

$$A[p][\langle u \rangle] = A$$

$$A[p][\gamma \uparrow] = A[\gamma][p]$$

$$A = A[p \uparrow][\langle \text{var } vz \rangle]$$

$$\text{app } (\text{lam } t) u = t[\langle u \rangle]$$

$$\frac{\text{Tm } \Delta ((\Pi A B)[\sigma])}{(\text{lam } t)[\sigma]} = \frac{\text{Tm } \Delta (\Pi (A[\sigma]) (B[\sigma \uparrow]))}{\text{lam } (t[\sigma \uparrow])}$$

$$\frac{\text{Tm } \Delta (B[\langle u \rangle][\sigma])}{(\text{app } t u)[\sigma]} = \frac{\text{Tm } \Delta (B[\sigma \uparrow][\langle u[\sigma] \rangle])}{\text{app } (t[\sigma])(u[\sigma])}$$

$$(\text{var } x)[p] = \text{var } (\text{vs } x)$$

$$(\text{var } vz)[\langle u \rangle] = u$$

$$(\text{var } (\text{vs } x))[\langle u \rangle] = \text{var } x$$

$$(\text{var } vz)[\gamma \uparrow] = \text{var } vz$$

$$(\text{var } (\text{vs } x))[\gamma \uparrow] = (\text{var } x)[\gamma][p]$$

$$t = \text{lam } (\text{app } (t[p])(\text{var } vz))$$

Details

We have a generalized algebraic theory

The syntax is the initial algebra

Other presentations:

- Thomas Ehrhard's thesis (1988)
- category with families (Dybjer 1995)
- natural model (Awodey 2016)
- contextual category (Brunerie 2019)
- B-system and C-system (Ahrens et al. 2023)
- Coquand also discovered single substitution calculus independently

CwF models are not equivalent to ours, but the syntaxes are isomorphic

Formalization in Agda with postulated QIIT and SProp

Summary

Minimalistic definition of type theory

Any SOGAT has a single substitution GAT presentation (Kaposi & Xie: FSCD 2024)

Need to derive parallel substitutions to prove normalization

Future work: coherent syntax of type theory?

$$(e : B[\langle u \rangle][\sigma] = B[\sigma \uparrow][\langle u[\sigma] \rangle]) \rightarrow (\text{app } t u)[\sigma] =_e \text{app } (t[\sigma]) (u[\sigma])$$